# High-dimensional prediction: Some computational challenges

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Baker-Kingland Lecture Predictive Inference and its Applications

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# The Data Explosion



- Every day: 2.5 billion gigabytes of data created
- $\circ~$  Last two years: creation of 90% of the world's data
- Data stored grows 4X faster than world economy (source: Mayer-Schonberger)

(source: IBM)

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- Interesting trade-offs between computational and statistical efficiency.

#### Today's talk: Two vignettes

- $\S1$  Data sketches: randomized dimensionality reduction
- $\S 2$  Early stopping of iterative algorithms for prediction

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- Project data into subspace, and solve reduced dimension problem.



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#### Randomized projection is a general purpose tool:

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- Project data into subspace, and solve reduced dimension problem.



Widely studied and used:

- Johnson & Lindenstrauss (1984): in Banach/Hilbert space geometry
- various surveys and books: Vempala, 2004; Mahoney et al., 2011 Cormode et al., 2012.

DATA

OPTIMIZER















#### Randomized projection for constrained least-squares

- Given data matrix  $A \in \mathbb{R}^{n \times d}$ , and response vector  $y \in \mathbb{R}^n$
- $\circ\,$  Least-squares over convex constraint set  $\mathcal{C}\subseteq \mathbb{R}^d\colon$

$$x_{\text{LS}} = \arg\min_{x \in \mathcal{C}} \underbrace{\|Ax - y\|_2^2}_{f(Ax)}$$

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• Randomized approximation:

(Sarlos, 2006, Mahoney et al., 2011)

$$\widehat{x} = \arg\min_{x \in \mathcal{C}} \|S(Ax - y)\|_2^2$$

• Random projection matrix  $S \in \mathbb{R}^{m \times n}$ 



# **Application to Netflix data**

NETFLIX				Your Account & Help Movies, TV shows, actors, directors, genres
Watch Instantly	Browse DVDs	Your Queue	Movies You'll 🖤	
Con Add mo Spider-Mar	gratulatio	ns! Movies of a source of the	we think You ve seen for even bett The Rundown	will 🖤 er suggestions. Bad Boys II



Las Vegas: Season 2 (6-Disc Series)















#### Netflix data set

- $\circ~2$  million  $\times$  17000 matrix A of ratings (users  $\times$  movies)
- Predict the ratings of a particular movie
- $\circ\,$  Least-squares regression with  $\ell_2$  regularization

$$\min_{x \in \mathbb{R}^{17000}} \left\{ \|Ax - y\|_2^2 + \lambda \|x\|_2^2 \right\}$$

• Partition into test and training sets, solve for all values of  $\lambda \in \{1, 2, ..., 100\}.$ 

## Sketching for Netflix movie database

original data set: 2 million × 17000 matrix A of ratings (users × movies)
o perform sketching (randomized dimensionality reduction)



#### Key fact:

Sketching to dimension 2000 is enough! Sketch is  $2000 \times 17000$ : a few Megabytes.

### Fitting the full regularization path

(Pilanci & W., 2016, J. Machine Learning Research)



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#### Iterative sketching for general data-based objectives

**Goal:** Minimize g(x) = f(Ax) over convex set  $\mathcal{C} \subseteq \mathbb{R}^d$ :

 $x_{\scriptscriptstyle \rm opt} = \arg\min_{x\in\mathcal{C}}g(x), \quad \text{where } g: \mathbb{R}^d \to \mathbb{R} \text{ is twice-differentiable}.$ 

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**Ordinary Newton steps:** 

$$x^{t+1} = \arg\min_{x\in\mathcal{C}} \left\{ \frac{1}{2} \|\nabla^2 g(x^t)^{1/2} (x - x^t)\|_2^2 + \langle \nabla g(x^t), x - x^t \rangle \right\},$$

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Sketched Newton steps: Using random sketch matrix  $S^t$ :

$$\tilde{x}^{t+1} = \arg\min_{x\in\mathcal{C}} \left\{ \frac{1}{2} \| S^t \nabla^2 g(x^t)^{1/2} (x - \tilde{x}^t) \|_2^2 + \langle \nabla g(\tilde{x}^t), \, x - \tilde{x}^t \rangle \right\}.$$

Cost per step:  $\widetilde{\mathcal{O}}(nd)$  in unconstrained case.

Run algorithm with sketch dimension  $m \asymp d$  on a self-concordant function g(x) = f(Ax), and data matrix  $A \in \mathbb{R}^{n \times d}$  with  $n \gg d$ .

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#### Theorem (Pilanci & W, SIAM J. Opt, 2017)

With probability at least  $1 - c_0 e^{-c_1 m}$ , number of iterations required for  $\epsilon$  accuracy is less than

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Dependence on sample size n, dimension d; conditioning  $\kappa$ ; and tolerance  $\epsilon$ 

Algorithm	Computational cost
Gradient Descent	$\mathcal{O}(\kappa  n  d  \log(1/\epsilon))$
Acc. gradient Descent	$\mathcal{O}(\sqrt{\kappa} nd \log(1/\epsilon))$
Newton's Method	$\mathcal{O}(nd^2 \log \log(1/\epsilon))$
Newton Sketch	$\widetilde{\mathcal{O}}(nd\log(1/\epsilon))$

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Note: Dependence on condition number  $\kappa$  unavoidable among 1st-order methods (Nesterov, 2004)

### Logistic regression: uncorrelated features



Sample size n = 500,000 with d = 5,000 features

### Logistic regression: correlated features



Sample size n = 500,000 with d = 5,000 features

## $\S.$ 2. Non-parametric regression via boosting

Non-parametric regression problem: approximate the regression function  $f^*(x) = \mathbb{E}[Y \mid X = x]$  based on samples  $\{(x_i, y_i)\}_{i=1}^n$ .

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 $\begin{array}{ll} \text{Empirical loss function} & \text{Function class } \mathscr{F} \\ \mathcal{L}_n : \mathscr{F} \to \mathbb{R} & \text{Norm } \| \parallel_{\mathscr{F}} \end{array}$ 

Given step sizes  $\alpha^t > 0$ :

$$f^{t+1} = f^t - \alpha^t g^t \qquad \text{where } g^t = \arg \max_{\|g\|_{\mathscr{F}} \leq 1} \langle g, \, \nabla \mathcal{L}_n(f^t) \rangle$$

[Freund & Schapire, 1997; Mason et al., 1999; Friedman et al., 2000]

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**Example:**  $L^2$ -Boosting with  $\mathcal{L}_n(f) = \frac{1}{2n} \sum_{i=1}^n [y_i - f(x_i)]^2$ . Gradient boosting update takes form

$$g^{t} = \arg \max_{\|g\|_{\mathscr{F}} \leq 1} \left\{ \frac{1}{n} \sum_{i=1}^{n} g(x_{i}) \underbrace{\left[ f^{t}(x_{i}) - y_{i} \right]}_{\text{Current residual}} \right\}$$

#### True function and noisy observations



#### Residuals: 1 rounds of Gauss kernel boosting





#### Function fit: 20 rounds of Gauss kernel boosting



#### Function fit: 40 rounds of Gauss kernel boosting



## Visualization of over-fitting



I shall not today attempt further to define the kinds of material I understand to be embraced within that shorthand description "overfitting", and perhaps I could never succeed in intelligibly doing so. But I know it when I see it.

Paraphrased from US Supreme Court Justice Potter Stewart, 1964

#### How to stop at the "right time"?



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#### **Desiderata:**

- a data-dependent stopping rule  $\{x_i, y_i\}_{i=1}^n \mapsto T \in \{1, 2, \dots, \}$
- $\circ$  function estimate at iteration t has optimal mean-squared error

$$||f^T - f^*||_n^2 \approx \min_{k=1,2,\dots} ||f^k - f^*||_n^2$$

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#### Some past work:

- o optimal but not data-dependent rule for spline kernel  $L^2\mbox{-boosting:}$  Bühlmann & Yu, 2003
- $\circ\,$  some results on  $L^2\mbox{-boosting}$  with spline kernels: Caponetto et al., 2007, Roscasco et al., 2009
- $\circ\,$ optimal rates for  $L^2$ -boosting, data-dependent any reproducing kernel: Raskutti, W. & Y., 2013

#### Kernel eigenvalues control "richness"



## Statistical error determined by fixed point equation



Fixed point equation:

$$\frac{1}{\sqrt{n}}\sqrt{\sum_{j=1}^{n}\min\left\{1,\ \frac{\widehat{\mu}_{j}^{2}}{\delta^{2}}\right\}} = \frac{\delta}{\sigma}$$

where

- $\circ \ \{\widehat{\mu}_j\}_{j=1}^n \text{ are eigenvalues} \\ \text{ of kernel matrix}$
- $\circ \ \sigma > 0$  is noise level

## Required number of iterations scales inversely...

Kernel	Stat. error	Number of iterations
Polynomial	$\frac{D}{n}$	$\frac{n}{D}$
(Degree $D$ )	10	D
Gaussian	$\frac{\sqrt{\log n}}{n}$	$\frac{n}{\sqrt{\log n}}$
		Viogn
First-order spline	$\left(\frac{1}{n}\right)^{\frac{2}{3}}$	$n^{2/3}$
-	(11)	
Second-order spline	$\left(\frac{1}{n}\right)^{\frac{4}{5}}$	$n^{4/5}$

### Minimax-optimal bounds for kernel boosting

Kernel boosting sequence:

$$f^{t+1} = f^t - \alpha g^t$$
 where  $g^t = \arg \max_{\|g\| \not \gg \le 1} \langle g, \nabla \mathcal{L}_n(f^t) \rangle$ 

#### Theorem (Wei, Yang & W., 2017)

For any kernel class  $\mathscr{F}$ , any (m, L)-regular loss function and any step size  $\alpha \in (0, \frac{m}{L}]$ , and any iterate  $t = 1, \ldots, \lfloor \frac{1}{\delta_{\alpha}^2} \rfloor$ :



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Statistical error determined by fixed point equation:

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where  $\{\widehat{\mu}_j\}_{j=1}^n$  are eigenvalues of kernel matrix, and  $\sigma > 0$  is noise level.

#### LogitBoost: Error on linear scale



### LogitBoost: Error on logarithmic scale



# Summary

Many challenges lie at the interface of statistics and optimization:

- data sketches for randomized dimension reduction
- regularization via early stopping of iterative algorithms for optimization

Some papers:

- Pilanci & W. (2016). Iterative Hessian sketch: Fast and accurate solution approximation for constrained least squares J. Machine Learning Research.
- Pilanci & W. (2017). Newton sketch: A linear-time optimization algorithm with linear-quadratic convergence. *SIAM Journal on Optimization*.
- Wei, Yang & W. (2017). Early stopping for kernel boosting algorithms. Arxiv pre-print.

#### Fast Johnson-Lindenstrauss sketch

**Step 1:** Choose some fixed orthonormal matrix  $H \in \mathbb{R}^{n \times n}$ . Example: Hadamard matrices

$$H_2 = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1\\ 1 & -1 \end{bmatrix} \qquad H_{2^t} = \underbrace{H_2 \otimes H_2 \otimes \cdots \otimes H_2}_{V_{\text{remarkup product } t \text{ time}}$$

Kronecker product t times

(E.g., Ailon & Liberty, 2010)

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#### Step 2:

- (A) Multiply data vector y with a diagonal matrix of random signs  $\{-1,+1\}$
- (B) Choose *m* rows of *H* to form sub-sampled matrix  $\widetilde{H} \in \mathbb{R}^{m \times n}$
- (C) Requires  $\mathcal{O}(n \log m)$  time to compute sketched vector  $Sy = \widetilde{H} Dy$ .

(E.g., Ailon & Liberty, 2010)

### Tools for sharp analysis



#### Localized Gaussian complexity

How much can you align with i.i.d. noise sequence  $\{w_i\}_{i=1}^n \sim N(0,1)$ ?

$$\mathscr{G}_n(\delta, r; \mathscr{F}) = \mathbb{E}_w \sup_{\substack{\|f\|_{\mathscr{F}} \leq r \\ \|f - f^*\| \leq \delta}} \left| \frac{1}{n} \sum_{i=1}^n w_i \big( f(x_i) - f^*(x_i) \big) \right|$$

(e.g., van de Geer, 2000; Bartlett et al., 2005; Koltchinski, 2006)